Binary Decisions Diagrams for Algebraic Attacks

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May 7, 2014 Winter School in Information Security Finse, Norway

Current State



Figure : Development of Algebraic Attack

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Binary Decisions Diagram (BDD)

$$f(x_1, x_2, x_3) = x_1 x_3 + x_1 + x_2 + x_3 + 1$$



Figure : Binary decision diagram for f function

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BDD in Cryptology



Figure : Example of a BDD with four levels

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S-box Representation Using BDD

$$S-box = \{5, C, 8, F, 9, 7, 2, B, 6, A, 0, D, E, 4, 3, 1\}$$



Binary decision diagram (BDD)-based cryptanalysis of

- A5/1 (GSM keystream generator)
- E0 (Bluetooth keystream generator)
- Trivium (eSTREAM Portfolio, Profile 2)
- Grain (eSTREAM Portfolio, Profile 2)

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N.T. Courtois, G.V. Bard [1]

The 6-round DES (with 20 fixed key bits) was attacked by algebraic attack in several minutes with the help of conversion to SAT and applying MiniSat 2.0.

E. Kleiman, [2]

The MiniAES (16-bit version) was attacked by XL and XSL methods.

"This results in a large sparse system of linear equations over the field GF(2) with an unknown number of extraneous solutions that need to be weeded out."

New Results of AA via BDD Representation

DES

- Our best result is finding the key of 6-round DES using 8 chosen plaintext/ciphertext pairs without fixing or guessing any variables.
- The average complexity is $2^{20.571}$ nodes, which is equivalent to ~ 1 minute on MacBook Air 2013 with 8GB RAM.

MiniAES

10-round MiniAES was totally broken via the BDD method using 1 known plaintext/ciphertext pair on regular PC. The average memory complexity is $2^{24.961}$ nodes.

Table : Complexities for solving reduced-round DES-systems. Each cell shows the minimum, **average** and maximum complexity observed over 100 instances.

# texts rounds	1	2	3	4	5	6	7	8
4	$2^{22.651}$ $2^{22.715}$	2 ^{10.800} 2 ^{14.506}	2 ^{9.281} 2 ^{10.606}	2 ^{9.585} 2 ^{10.257}	2 ^{9.748} 2 ^{9.805}	2 ^{9.976} 2 ^{10.070}	2 ^{10.103} 2 ^{10.203}	2 ^{10.283} 2 ^{10.381}
	222.110	210.479	213.000	212.029	25.052	210.412	210.578	210.440
5		$2^{19.472}$	$2^{13.031}$	$2^{11.440}$	212.120	$2^{12.209}$	$2^{12.565}$	$2^{12.749}$
		$2^{23.805}$	$2^{10.400}$ $2^{19.329}$	$2^{15.618}$ $2^{15.618}$	$2^{16.633}$ $2^{16.633}$	$2^{16.758}$	$2^{14.410}$ $2^{16.882}$	$2^{17.414}$ $2^{17.414}$
6						$2^{24.506}$	$2^{22.206}$	$2^{19.932}$
						$2^{24.929}$	$2^{22.779}$	$2^{20.571}$
						$2^{25.352}$	$2^{24.324}$	$2^{21.915}$

Open Problems and Further Development

- Development of general methodology and justification of theoretical bounds:
 - Does there exist a generic algorithm giving an order of BDDs that yield low complexity when applying linear absorption?
 - Is it possible to analytically estimate the complexity of solving a BDD system of equations, or do we have to actually run the solver to find out?
 - Which ciphers are most vulnerable against this type of algebraic attacks?
- More block ciphers, stream ciphers and hash functions can be attacked

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- Courtois, N.T., Bard, G.V., Algebraic cryptanalysis of the Data Encryption Standard, Cryptography and Coding, LNCS 4887, pp. 152–169, Springer (2007).
- Kleiman, E., High Performance Computing techniques for attacking reduced version of AES using XL and XSL methods, Graduate Theses and Dissertations (2010) http://lib.dr.iastate.edu/etd/11473.